REMARKS

The application includes claims 1-17 and 26-38 prior to entering this amendment.

The applicants amend claims 1, 3-5, 7-15, 26, 29, 32, 33 and 36.

The applicants cancel claims 2, 16-17.

The applicants add claims 39-41.

The application remains with claims 1, 3-15, 26-28 and 30-41 after entering this amendment.

The above amendments are made without prejudice or disclaimer. The amendments are made to more clearly delineate intended subject matter. Accordingly, the applicants do not intend to surrender claimed subject matter by submission of the above amendments and do not add new matter. The applicants respectfully requests reconsideration of the above referenced patent application in view of the following remarks.

Claim Rejections - 35 U.S.C. § 103

The office action rejected claims 1-17 and 26-38 under 35 U.S.C. § 103(a) as being unpatentable over Dykeman *et al.* (U.S. Patent 7,177,951) in view of Rajsic *et al.* (U.S. Patent 7,283,467). Claims 1, 3-15 and 26-38 are currently pending. Applicants submit that claims 1, 3-15 and 26-38 are allowable over Dykeman in view of Rajsic. Therefore, Applicants respectfully traverse the claim rejections for the reasons explained herein.

In the current application applicants claim a crankback method wherein upon detecting a call failure *inside* a succeeding peer group, a succeeding end crankback specifies a blocked interface on a first link between preceding and succeeding peer groups. The succeeding end crankback causes the first peer group to send a second connection request using a second link to the second peer group that avoids the call failure.

With respect to amended claim 1, applicants claim:

> transmitting a succeeding end crankback rather than a next higher level crankback from the succeeding peer group to the preceding peer group wherein transmitting is based at least in part on the determination of whether multiple nodes in the succeeding peer group have connectivity to the preceding peer group or whether multiple nodes of the preceding peer group have connectivity to the first node of the succeeding peer group, or combinations thereof *(Emphasis added)*

end crankback rather than a next higher level crankback. *See* the Office Action, page 2. Applicants agree. Therefore, the remainder of the remarks will address the deficiencies of Rajsic with respect to crankback procedure.

The office action cites Rajsic for the feature, "transmitting a succeeding end crankback rather than a next higher level crankback from the succeeding peer group to the preceding peer group" of claim 1. However, the crankback procedure of Rajsic is not the same as the crankback procedure of the current application as claimed in amended claim 1.

In Rajsic, a preceding peer group and a succeeding peer group are connected via a first link. When a call fails in the succeeding peer group, a succeeding end crankback is sent without a *determination of whether multiple nodes in the succeeding peer group have connectivity to the preceding peer group* or *whether multiple nodes of the preceding peer group have connectivity to the first node of the succeeding peer group* as is disclosed in claim 1. In other words, Rajsic's method does not check to see if alternate links for connecting the preceding peer group and the succeeding peer group are available before sending the succeeding end crankback. Thus, the procedure of Rajsic is random and less efficient than the procedure disclosed in the current application because in Rajsic a node *may try* parallel trunks to set-up a second link between the preceding peer group and the succeeding peer group unsuccessfully and *then* revert to conventional crankback procedures. See Rajsic, col. 5, lines 65-67. Further, the succeeding end crankback of Rajsic has to be converted to a conventional crankback if sent erroneously, whereas the succeeding end crankback as claimed in claim 1 of the current application is not sent unless it is determined that a second link is possible. Likewise, in the current disclosure, a determination that there are potential second links is made prior to substituting the succeeding end crankback thus determining that a second link is possible.

Additionally, the modified crankback procedure in Rajsic is limited with respect to the disclosed crankback procedure of the current application. Specifically, in the modified crankback procedure of Rajsic, only the node of the preceding peer group that was *originally linked* to the succeeding peer group can try setting up a second link to the succeeding peer group. Accordingly, if the *originally linked node* of the preceding peer group does not have other parallel trunks connected to the succeeding peer group, the entire succeeding peer group is treated as blocked (as in conventional crankback procedures) because this scenario converts the modified crankback of Rajsic to a conventional crankback method. See Rajsic, Fig. 3, col. 5, line

treated as blocked (as in conventional crankback procedures) because this scenario converts the modified crankback of Rajsic to a conventional crankback method. See Rajsic, Fig. 3, col. 5, line 65-col. 6, line 2 and Fig. 4, lines 3-11. Thus, even if there are *other nodes* in the preceding peer group in Rajsic that are connected to the succeeding peer group the other nodes of the preceding peer group will not be used to set-up a second link and the entire succeeding peer group will be treated as blocked.

In contrast, amended claim 1 discloses:

> wherein the second link is between the first node of the preceding peer group and a second node of the succeeding peer group or a second node of the preceding peer group and a second node of the succeeding peer group *(Emphasis added)*

Thus, claim 1 discloses that alternate nodes of the preceding peer group may be used to set up a second link to the succeeding peer group thereby opening many more opportunities to access unused resources than may be used in a network implementing conventional crankback procedures or the modified crankback procedures of Rajsic.

Further, even if Rajsic did teach the features of claim 1 as is asserted in the office action (which it does not for the reasons explained above), Dykeman teaches away from the proposed modification because Dykeman does not even address crankback procedures other than to disclose that unnecessary crankback procedures may be avoided by implementation of the method of Dykeman for address management wherein a peer group leader can actively check address connectivity within a peer group and update its topology database rather than waiting for address information to be flooded. *See* Dykeman col. 3, lines 15-43. The method of only Dykeman avoids unnecessary crankback because the address information is kept current between standard address updating procedures. In Dykeman, any mention of crankback relates to conventional crankback procedures that do not involve setting up new links *between* different peer groups as is disclosed in claim 1. *See* Dykeman col. 10, lines 60-67 and col. 11, lines 1-12 and see Fig. 5.

Accordingly, claim 1 distinguishes from Dykeman and Rajsic and should be allowed.

Claims 3-8 depend from claim 1 and should also be allow.

Claims 9, 26 and 36 and the respective dependent claims distinguish from Dykeman and Rajsic on at least the same or similar basis and should be allowed.

New claims 39-41 contain similar features to the features of claim 1 and therefore distinguish from Dykeman and Rajsic for at least the same or similar reasons as claim 1 and should be allowed.

New Claims

New claims 39-41 have been added and are supported throughout the disclosure.

New claims 39-41 are allowable over Dykeman and Rajsic. For instance, claim 39 contains the feature "prior to sending a first type of crankback message to the source for the detected call disruption, determine whether one of the peer switches besides the destination peer switch has connectivity with the remote logical group... [and] if one of the peer switches besides the destination peer switch has connectivity with the remote logical group, do not send the first type of crankback message, and instead send a second different type of crankback message, wherein the second type of crankback message indicates communication disruption over the inter group link when in fact the communication disruption is over one of the intra group links, the second type of crankback message indicating an inter group link failure even when there is no inter group link failure, the second type of crankback message configured to cause the source to reroute the call using a different inter group link..." which is not disclosed in Dykeman or Rajsic for reasons similar to those discussed above. Therefore, new claim 39 should be allowed. Claims 40-41 depend from claim 39 and should also be allowed.

CONCLUSION

For the foregoing reasons, the applicants request reconsideration and allowance of all pending claims. The applicants encourage the examiner to telephone the undersigned if it appears that an interview would be helpful in advancing the case.

Customer No. 73552

Respectfully submitted,

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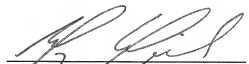

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